

# Structure and enumeration theorems for hereditary properties in finite relational languages

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## Logical 0-1 laws

Fix a finite first-order language  $\mathcal{L}$ . For each  $n$  suppose  $K(n)$  is a set of  $\mathcal{L}$ -structures with universe  $[n] = \{1, \dots, n\}$ . Set  $K = \bigcup_{n \in \mathbb{N}} K(n)$ .

### Definition

We say  $K$  has a 0-1 law if for every  $\mathcal{L}$ -sentence  $\phi$ ,

$$\mu(\phi) = \lim_{n \rightarrow \infty} \frac{|\{G \in K(n) : G \models \phi\}|}{|K(n)|}$$

is either 0 or 1.

The *almost sure theory* of  $K$  is  $T_{as}(K) = \{\phi \in \mathcal{L} : \mu(\phi) = 1\}$ .

If  $K$  has a 0-1 law, then  $T_{as}(K)$  is complete.

# Discrete Metric Spaces

This is joint work with D. Mubayi.

## Definition

Let  $r \geq 2$  be an integer.

- 1  $M_r(n)$  is the set of metric spaces with underlying set  $[n]$  and distances all in  $[r]$ .
- 2  $\mathcal{L}_r = \{d_1, \dots, d_r\}$  where each  $d_i$  is a binary relation symbol.

Every  $G \in M_r(n)$  is naturally an  $\mathcal{L}_r$ -structure:  
for all  $a, b \in G$ , interpret

$$G \models d_i(a, b) \Leftrightarrow d(a, b) = i \text{ in } G.$$

## Question

Does  $M_r := \bigcup_{n \in \mathbb{N}} M_r(n)$  have a 0-1 law?

## Answer for $r$ even

**Assume  $r \geq 2$  is even.** We now define a special subfamily of  $M_r(n)$ .

### Definition

Let  $C_r(n) = \{G \in M_r(n) : \text{for all } a \neq b \in [n], d(a, b) \in \{\frac{r}{2}, \dots, r\}\}$ .

### Theorem (Mubayi, T.)

$$\lim_{n \rightarrow \infty} \frac{|C_r(n)|}{|M_r(n)|} = 1.$$

Let  $C_r = \bigcup_{n \in \mathbb{N}} C_r(n)$ . A standard argument shows  $C_r$  has a 0-1 law.

### Corollary (Mubayi, T.)

$M_r$  has a 0-1 law and  $T_{as}(M_r) = T_{as}(C_r)$ .

Idea: Precise structure theorem + 0-1 law for  $C_r \Rightarrow$  new 0-1 law for  $M_r$ .

# How do we prove the precise structure theorem?

Key tool: approximate structure and enumeration.

## Definition

Given  $\delta > 0$  and two elements  $G, G' \in M_r(n)$ , we say  $G$  and  $G'$  are  $\delta$ -close if  $\left| \left\{ ab \in \binom{[n]}{2} : d^G(a, b) \neq d^{G'}(a, b) \right\} \right| \leq \delta \binom{n}{2}$ .

Let  $C_r^\delta(n) = \{G \in M_r(n) : G \text{ is } \delta\text{-close to an element of } C_r(n)\}$ .

## Theorem (Mubayi, T.)

*Structure: for all  $\delta > 0$ , there is  $\beta > 0$  such that for large  $n$ ,*

$$\frac{|M_r(n) \setminus C_r^\delta(n)|}{|M_r(n)|} \leq 2^{-\beta \binom{n}{2}}.$$

*Enumeration:  $|M_r(n)| = |C_r(n)|(1 + 2^{o(n^2)}) = \left(\frac{r}{2} + 1\right) \binom{n}{2} + o(n^2)$ .*

# Outline of 0-1 law

Approximate Structure and Enumeration + Ad-hoc arguments



Exact Structure and Enumeration ( $C_r(n)$  takes over)

+

0-1 law for less complicated family ( $C_r(n)$ )



0-1 law for complicated family ( $M_r(n)$ )

Examples where one can apply this strategy:

- 1  $K_\ell$ -free graphs (Kolaitis-Prömel-Rothschild)
- 2  $T_k$ -free digraphs (Kühn-Osthus-Townsend-Zhao) + (Koponen)
- 3 Triangle-free 3-uniform hypergraphs (Balogh-Mubayi) + (Koponen)
- 4 Discrete metric spaces (Mubayi-T.)

# Focus on Approximate Structure and Enumeration

Approximate structure and enumeration results \_\_\_\_\_.

- are important tools in proofs of certain 0-1 laws.
- are of independent interest in extremal combinatorics.
- have been proven for lots of combinatorial objects (e.g. graphs, digraphs, hypergraphs, colored hypergraphs, metric spaces).
- often have similar proofs using combination of:
  - 1 extremal results
  - 2 supersaturation results
  - 3 stability theorems
  - 4 graph removal lemmas
  - 5 regularity lemmas
  - 6 hypergraph containers theorem (Balogh-Morris-Samotij, Saxton-Thomason).

# Question

## Question

Is there a way to view these results (and their proofs) as examples of a general theorem (and its proof)?

Today: yes.

Main Ingredients:

- Hypergraph containers theorem (Balogh-Morris-Samotij, Saxton-Thomason).
- Triangle Removal for  $\mathcal{L}$ -structures (Aroskar-Cummings).
- Many combinatorics papers which have made the pattern of proof clear. Particularly recent work using the hypergraph containers theorem.

Remark: similar results were obtained independently by Falgas-Ravry, O'Connell, and Uzzell.



# Hereditary $\mathcal{L}$ -properties

Let  $\mathcal{L}$  be a finite relational language and  $\mathcal{H}$  a class of finite  $\mathcal{L}$ -structures.  $\mathcal{H}$  has the *hereditary property* if  $A \in \mathcal{H}$  and  $B \subseteq_{\mathcal{L}} A$  implies  $B \in \mathcal{H}$ .

## Definition

$\mathcal{H}$  is a *hereditary  $\mathcal{L}$ -property* if it has the hereditary property and is closed under isomorphism.

In the appropriate language, most of the results we want to generalize are for hereditary  $\mathcal{L}$ -properties.

Enumeration and structure of hereditary properties in the setting of graphs and other combinatorial structures have been studied in combinatorics.

# Setup

For the rest of the talk,

- $\mathcal{L}$  is a fixed, finite relational language.
- $r$  is the maximum arity of the relation symbols in  $\mathcal{L}$ .
- Assume  $\mathcal{H}$  is a hereditary  $\mathcal{L}$ -property.
- For each  $n$ ,  $\mathcal{H}_n$  is the set of elements in  $\mathcal{H}$  with domain  $[n]$ .
- For all  $n$ ,  $\mathcal{H}_n \neq \emptyset$ .

## Questions

- 1  $|\mathcal{H}_n| = ??$
- 2 What is the approximate structure of  $\bigcup_{n \in \mathbb{N}} \mathcal{H}_n$ ?

# $\mathcal{L}_{\mathcal{H}}$ -structures

## Definition

$S_r(\mathcal{H})$  is the set of complete, quantifier-free  $\mathcal{L}$ -types  $p(x_1, \dots, x_r)$  s.t. for each  $i \neq j$ ,  $x_i \neq x_j \in p(\bar{x})$  and  $p(\bar{x})$  is realized in some element of  $\mathcal{H}$ .

## Definition

$\mathcal{L}_{\mathcal{H}} = \{R_p(x_1, \dots, x_r) : p(x_1, \dots, x_r) \in S_r(\mathcal{H})\}$ .

Notation:  $V^{\underline{\ell}} = \{(a_1, \dots, a_{\ell}) \in V^{\ell} : a_i \neq a_j \text{ each } i \neq j\}$ .

## Definition

Let  $V$  be a set. An  $\mathcal{L}_{\mathcal{H}}$ -structure  $M$  with domain  $V$  is an  $\mathcal{L}_{\mathcal{H}}$ -*template* if

- For all  $\bar{a} \in V^{\underline{r}}$ , there is  $R_p(\bar{x}) \in \mathcal{L}_{\mathcal{H}}$  such that  $M \models R_p(\bar{a})$ .
- For all  $p, q \in S_r(\mathcal{H})$ , if  $p(x_1, \dots, x_r) = q(x_{\mu(1)}, \dots, x_{\mu(r)})$  for some permutation  $\mu$  of  $[r]$ , then for all  $(a_1, \dots, a_r) \in V^{\underline{r}}$ ,

$M \models R_p(a_1, \dots, a_r)$  if and only if  $M \models R_q(a_{\mu(1)}, \dots, a_{\mu(r)})$ .

## Example

$\mathcal{L} = \{E(x, y)\}$  and  $\mathcal{H}$  is the class of all finite triangle-free graphs.

$p(x, y) =$  the complete q.f. type containing  $E(x, y) \wedge E(y, x) \wedge x \neq y$ .

$q(x, y) =$  the complete q.f. type containing  $\neg E(x, y) \wedge \neg E(y, x) \wedge x \neq y$ .

Then  $S_r(\mathcal{H}) = \{p(x, y), q(x, y)\}$  and  $\mathcal{L}_{\mathcal{H}} = \{R_p(x, y), R_q(x, y)\}$ .

Draw Example.

# Subpatterns

Suppose  $M$  is an  $\mathcal{L}_{\mathcal{H}}$ -template with domain  $V$ .

## Definition

An  $\mathcal{L}$ -structure  $N$  is a *full subpattern* of  $M$  if  $\text{dom}(N) = V$  and for all  $\{a_1, \dots, a_r\} \in \binom{V}{r}$ ,

if  $p(x_1, \dots, x_r) = qftp^N(a_1, \dots, a_r)$  then  $M \models R_p(a_1, \dots, a_r)$ .

In this case, write  $N \trianglelefteq_p M$ .

Observe any  $\mathcal{L}$ -structure  $N$  with domain  $V$  is determined by  $qftp^N(a_1, \dots, a_r)$  for all  $\{a_1, \dots, a_r\} \in \binom{V}{r}$ .

Example.

# Extremal Structures

## Definition

An  $\mathcal{L}_{\mathcal{H}}$ -template  $M$  is called  $\mathcal{H}$ -random if  $N \trianglelefteq_p M$  implies  $N \in \mathcal{H}$ .

- $\mathcal{R}([n], \mathcal{H})$  is the set of  $\mathcal{H}$ -random  $\mathcal{L}_{\mathcal{H}}$ -templates with domain  $[n]$ .
- $sub(M) = |\{N : N \trianglelefteq_p M\}|$ .
- $ex(n, \mathcal{H}) = \max\{sub(M) : M \in \mathcal{R}([n], \mathcal{H})\}$ .

## Definition

$M \in \mathcal{R}([n], \mathcal{H})$  is *extremal* if  $sub(M) = ex(n, \mathcal{H})$ .

Observation:  $|\mathcal{H}_n| \geq ex(n, \mathcal{H})$ .

- $\mathcal{R}_{ex}([n], \mathcal{H})$  is the set of extremal  $M$  in  $\mathcal{R}([n], \mathcal{H})$ .

## Results: Enumeration

Recall  $\text{ex}(n, \mathcal{H}) = \max\{\text{sub}(M) : M \in \mathcal{R}([n], \mathcal{H})\}$ .

The *asymptotic density* of  $\mathcal{H}$  is  $\pi(\mathcal{H}) := \lim_{n \rightarrow \infty} \text{ex}(n, \mathcal{H})^{1/\binom{n}{r}}$ .

### Theorem (T.)

For all hereditary  $\mathcal{L}$ -properties  $\mathcal{H}$ ,  $\pi(\mathcal{H})$  exists.

### Theorem (T.)

If  $\mathcal{H}$  is a hereditary  $\mathcal{L}$ -property, then the following hold.

- 1 If  $\pi(\mathcal{H}) > 1$ , then  $|\mathcal{H}_n| = \pi(\mathcal{H})^{\binom{n}{r} + o(n^r)}$ .
- 2 If  $\pi(\mathcal{H}) \leq 1$ , then  $|\mathcal{H}_n| = 2^{o(n^r)}$ .

# Distance Between First-Order Structures

## Definition

Suppose  $\mathcal{L}_0$  is a finite relational language with maximum arity  $r$ . If  $M$  and  $N$  are finite  $\mathcal{L}_0$ -structures with domain  $V$ . Set

$$\text{dist}(M, N) = \frac{|\{A \in \binom{V}{r} : \text{qftp}^M(A) \neq \text{qftp}^N(A)\}|}{\binom{|V|}{r}}.$$

We say  $M$  and  $N$  are  $\delta$ -close if  $\text{dist}(M, N) \leq \delta$ .

This is basically the same as a definition of Aroskar-Cummings.



# Stability and Structure

## Definition

$\mathcal{H}$  has a *stability theorem* if for all  $\delta > 0$  there is  $\epsilon > 0$  such that for sufficiently large  $n$  the following holds. For all  $M \in \mathcal{R}([n], \mathcal{H})$

$$\text{sub}(M) \geq \text{ex}(n, \mathcal{H})^{1-\epsilon} \Rightarrow M \text{ is } \delta\text{-close to an element of } \mathcal{R}_{\text{ex}}([n], \mathcal{H}).$$

Let  $E(n, \mathcal{H}) = \{G \in \mathcal{H}_n : G \trianglelefteq_p M, \text{ some } M \in \mathcal{R}_{\text{ex}}([n], \mathcal{H})\}$ .

Let  $E^\delta(n, \mathcal{H}) = \{G \in \mathcal{H}_n : G \text{ is } \delta\text{-close to some } G' \in E(n, \mathcal{H})\}$ .

## Theorem (T.)

Suppose  $\pi(\mathcal{H}) > 1$  and  $\mathcal{H}$  has a stability theorem. Then for all  $\delta > 0$ , there is a  $\beta > 0$  such that for sufficiently large  $n$ ,

$$\frac{|\mathcal{H}_n \setminus E^\delta(n, \mathcal{H})|}{|\mathcal{H}_n|} \leq 2^{-\beta \binom{n}{r}}.$$

## Back to Metric Spaces

Given  $r \geq 2$ , let  $\mathcal{L} = \{d_1, \dots, d_r\}$ , and let  $\mathcal{M}_r$  be the hereditary  $\mathcal{L}$ -property obtained by closing  $\bigcup_{n \in \mathbb{N}} M_r(n)$  under isomorphism.

### Fact

If  $r$  is even then the following hold.

- $E(n, \mathcal{M}_r) = C_r(n)$ , so for all  $\delta > 0$ ,  $E^\delta(n, \mathcal{M}_r) = C_r^\delta(n)$ .
- $\pi(\mathcal{M}_r) = \left(\frac{r}{2} + 1\right)$ .

Consequently, our general counting theorem implies the following.

### Theorem

When  $r \geq 2$  is even,  $|M_r(n)| = \pi(\mathcal{M}_r) \binom{n}{2} + o(n^2) = \left(\frac{r}{2} + 1\right) \binom{n}{2} + o(n^2)$ .

# Back to Metric Spaces

## Fact

If  $r \geq 2$  is even then  $\mathcal{M}_r$  has a stability theorem.

Consequently, our general structure theorem implies the following.

## Theorem

Let  $r \geq 2$  be even. For all  $\delta > 0$ , there is  $\beta > 0$  such that for large  $n$ ,

$$\frac{|(\mathcal{M}_r)_n \setminus E^\delta(n, \mathcal{M}_r)|}{|(\mathcal{M}_r)_n|} = \frac{|M_r(n) \setminus C_r^\delta(n)|}{|M_r(n)|} \leq 2^{-\beta \binom{n}{2}}.$$

## Back to Metric Spaces

Recall that when  $r$  is even,  $\bigcup_{n \in \mathbb{N}} M_r(n)$  has a 0-1 law (Mubayi-T.).

### Conjecture (Mubayi-T.)

When  $r$  is odd,  $\bigcup_{n \in \mathbb{N}} M_r(n)$  does not have a 0-1 law.

Question: What is the difference between the even and odd cases?

### Theorem (Mubayi-T., T.)

*When  $r \geq 2$  is even,  $\mathcal{M}_r$  has a stability theorem. When  $r$  is odd,  $\mathcal{M}_r$  does not have a stability theorem.*

# A Conjecture

This naturally leads to the following conjecture.

Assume  $\mathcal{L}$  is a finite relational language.

## Definition

Suppose  $\mathcal{H}$  is a collection of finite  $\mathcal{L}$ -structures such that  $\mathcal{H}_n \neq \emptyset$  for each  $n$ . Then  $\mathcal{H}$  has a 0-1 law if  $\bigcup_{n \in \mathbb{N}} \mathcal{H}_n$  has a 0-1 law.

## Conjecture (T.)

If  $\mathcal{H}$  is a hereditary  $\mathcal{L}$ -property with  $\pi(\mathcal{H}) > 1$  and a 0-1 law, then  $\mathcal{H}$  has a stability theorem.



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